

CSE 390B, Winter 2023

Building Academic Success Through Bottom-Up Computing

Professional Networking & Compiler Phases

Professional Networking in College, Exploring the Compiler
Phases, Project 7 Overview

Lecture Outline

- ❖ **Professional Networking in College**
 - **Benefits of Building Connections, Networking Strategies**
- ❖ Exploring the Compiler Phases
 - Scanner: Process of Tokenizing an Input File
 - Parser: Making Meaning From Tokens Through ASTs
 - Type Checking, Optimization, and Code Generation
- ❖ Project 7 Overview
 - Midterm Corrections, Professor Meeting Report

Benefits of Building Connections

- ❖ Reaching out to your professors, TAs, and peers can be a great way to discover opportunities
- ❖ Taking the time to connect with these people can open several doors and leverage your potential
- ❖ Excellent opportunity for new perspectives and ideas for those who have been in your shoes before
- ❖ Connecting with others helps you find inspiration and build your knowledge and experience

Strategies for Networking

- ❖ Get involved in communities on campus (e.g., RSOs, TAing, research, part-time campus job)
- ❖ Invest in building relationships with people and developing a presence in their lives
- ❖ Take time to reflect on how others can support you by bringing to them your interests and questions
- ❖ Not all networking efforts will be well-received, but don't be afraid to just go for it

Discussion on Building Connections

Take some time to think about these questions:

- ❖ In what ways do you already connection with others on a regular basis? How else can you build your connections?
- ❖ How can you benefit from building your community of people you can network with?
- ❖ What would you share with someone you recently made a connection with?

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The Compiler: Goal

```
public int fact(int n) {  
    if (n == 0) {  
        return 1;  
    } else {  
        return n * fact(n - 1);  
    }  
}
```

High-Level Language

Theory Definition: a string, from the set of strings making up a language

Practical Definition: a file containing a bunch of characters

```
(fact)  
@R0  
M=M+1  
@R1  
D=A  
@ifbranch  
D;JEQ
```

Assembly Language



Compiler



The Compiler: Implementation

```
public int fact(int n) {  
    if (n == 0) {  
        return 1;  
    } else {  
        return n * fact(n - 1);  
    }  
}
```

High-Level Language

```
(fact)  
    @R0  
    M=M+1  
    @R1  
    D=A  
    @ifbranch  
    D;JEQ
```

Assembly Language

Scanner

Parser

Type
Checker

Optimizer

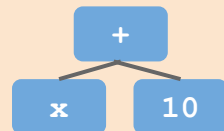
Code
Generator

Break string into
discrete **tokens**:

IF (ID(n)

== NUM(0) etc.

Arrange tokens into
syntax tree:



Verify the
syntax tree is
**semantically
correct**

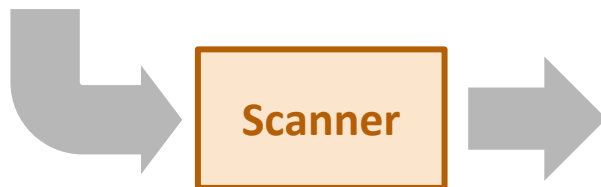
Rearrange the
code to be
more efficient

Convert the syntax
tree to the **target
language**

The Scanner

```
function void main() {  
  var int a, bar;  
  let bar=10; // init  
}
```

Jack



FUNCTION VOID ID (main)

LPAREN RPAREN LCURLY VAR

INT ID (a) COMMA ID (bar)

SEMICOLON LET ID (bar)

EQUALS NUM (10) SEMICOLON

RCURLY

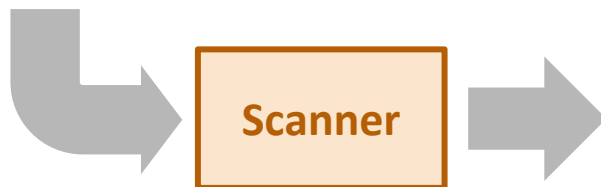
Token Stream

- ❖ Reads a giant string, breaks down into tokens
 - Each token has a type: what role does this token play?
 - E.g., **LCURLY** is a type representing an occurrence of “{”
 - What types do we care about? The “building blocks” of our programming language:
 - Keywords (e.g., **FUNCTION**), operators (e.g., **EQUALS**), and punctuation (e.g., **SEMICOLON** or **COMMA**)

The Scanner

```
function void main() {  
  var int a, bar:  
  let bar=10; // init  
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```

Jack



FUNCTION VOID ID (main)

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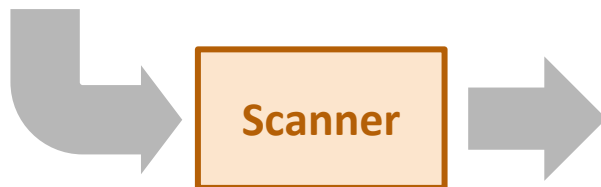
Token Stream

- ❖ In addition to a type, some tokens carry a value:
 - Identifiers (e.g., `ID (a)`)
 - Numbers (e.g., `NUM (10)`)
- ❖ Scanner should present a *clean* token stream
 - No whitespace or comments: the rest of the compiler only wants to consider things that change program meaning

The Scanner: How?

```
function void main() {  
  var int a, bar;  
  let bar=10; // init  
}
```

Jack



FUNCTION VOID ID (main)

LPAREN RPAREN LCURLY VAR

INT ID (a) COMMA ID (bar)

SEMICOLON LET ID (bar)

EQUALS NUM (10) SEMICOLON

RCURLY

Token Stream

- ❖ What if we split the input program on whitespace, and match each segment to a token type? (E.g., “{” → LCURLY)
- ❖ Tempting, but we would end up with “a,” “bar;” “bar=10;”
 - Whitespace is tricky: generally, we want to ignore it, but we can’t count on it being there

The Scanner: How?

curr



```
; let bar=10;
```

Jack

Accumulated: ;

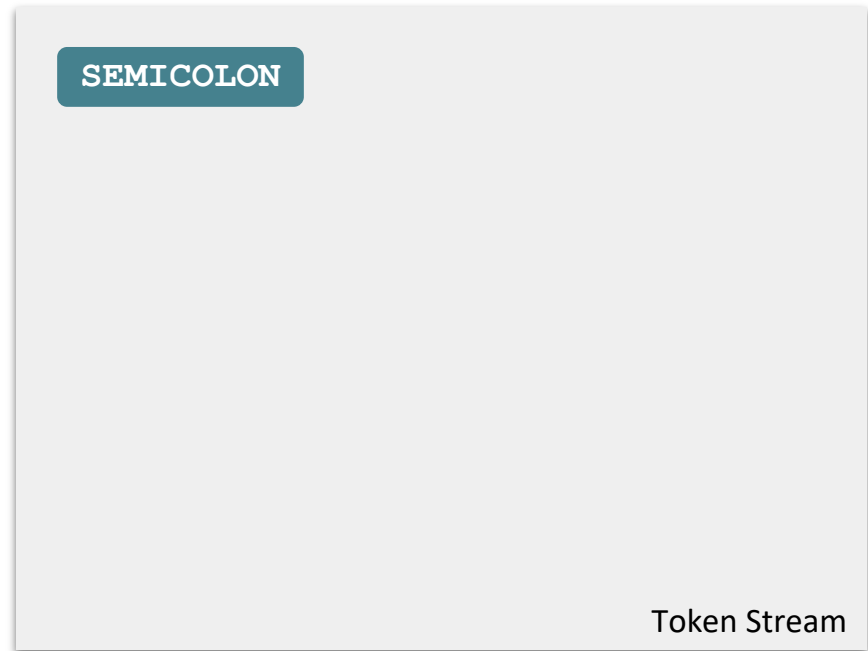
Token Stream

- ❖ How can we take a line of code in Jack and convert this into a token stream?
 - Keep cursor on current char
 - Break off a token when we complete one
 - If the next char could be part of this token, accumulate it

The Scanner: How?



Accumulated:

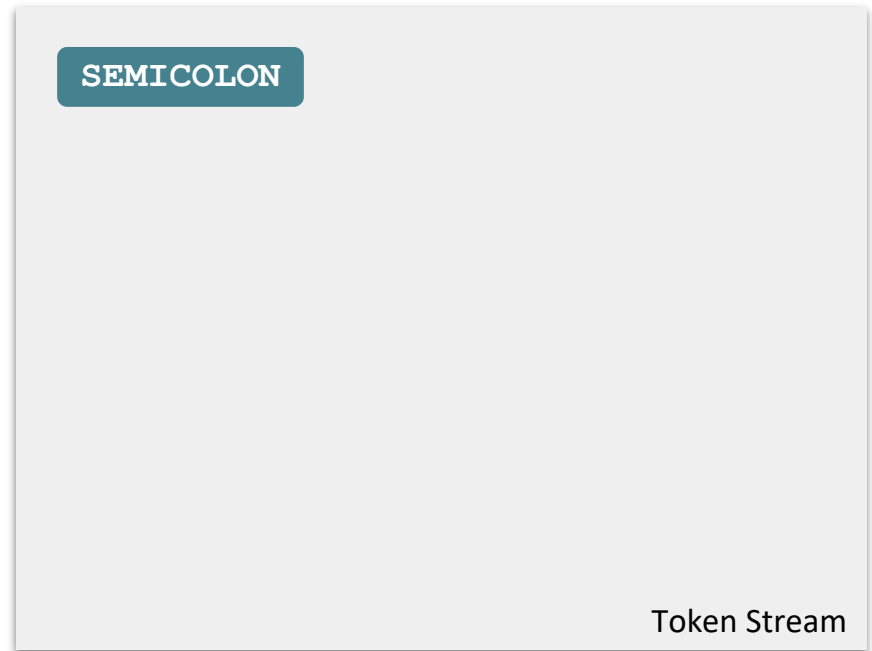


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The Scanner: How?



Accumulated: 1

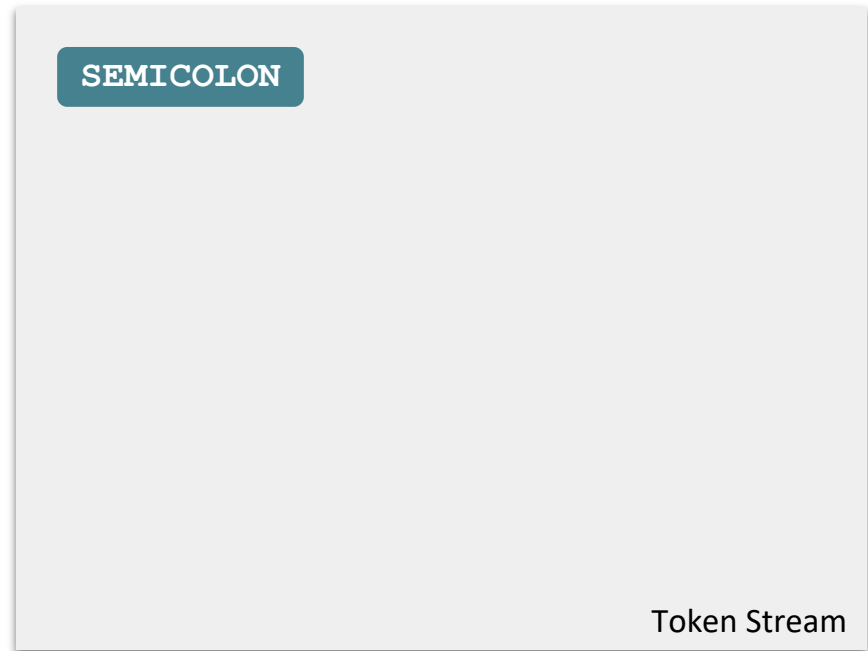


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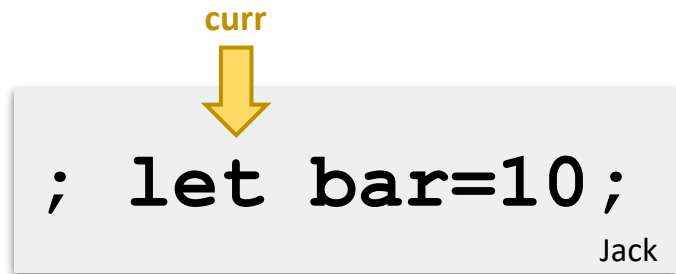


Accumulated: `le`

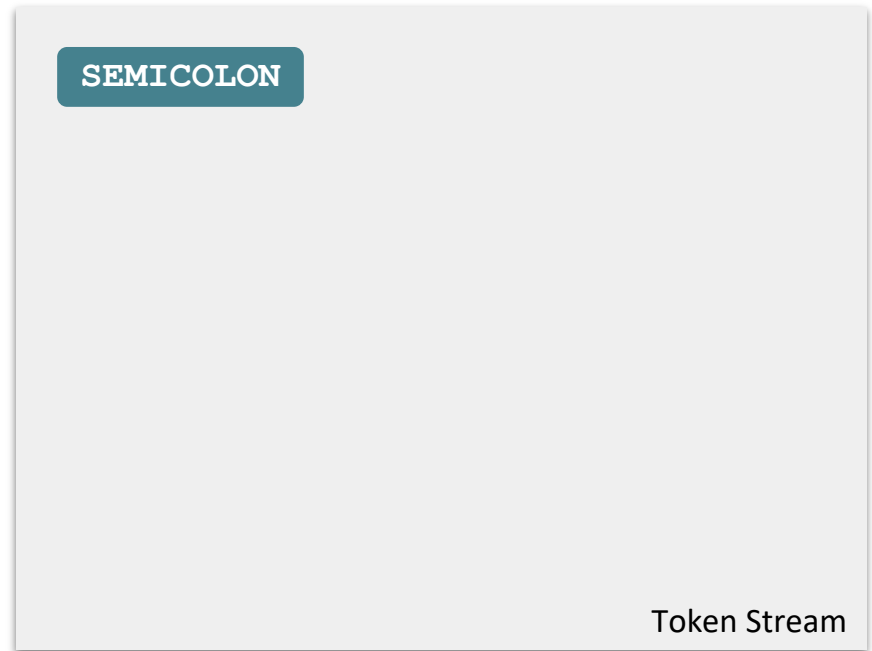


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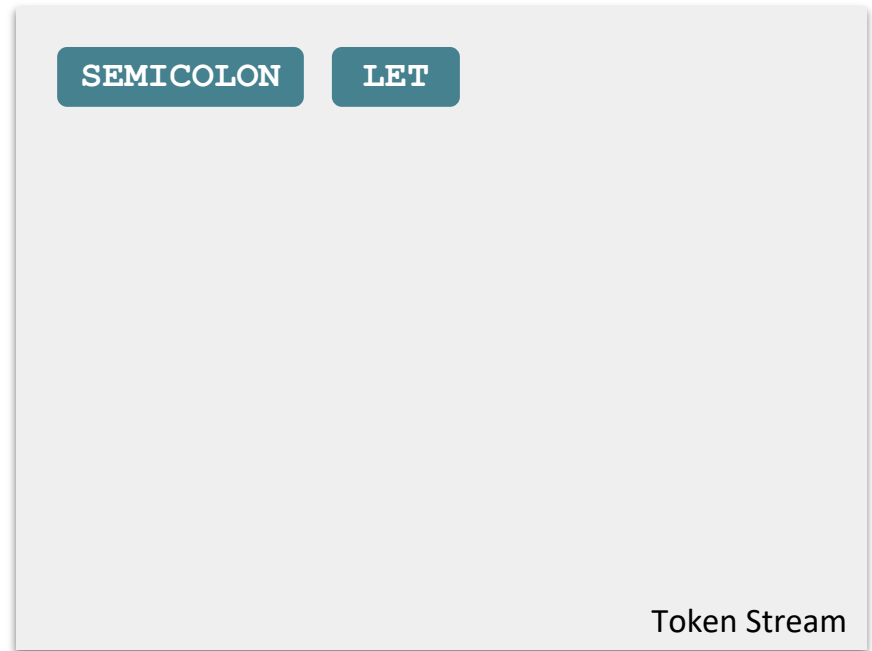
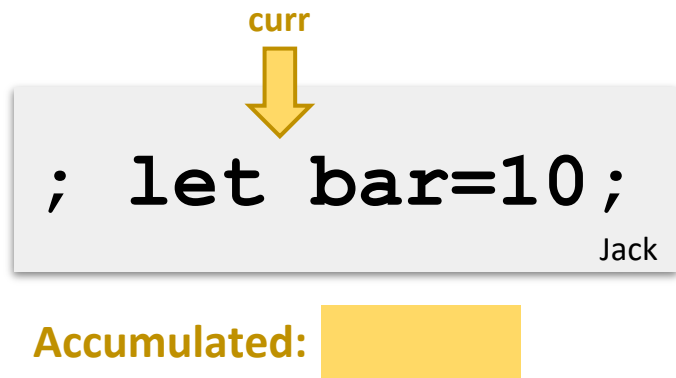


Accumulated: let



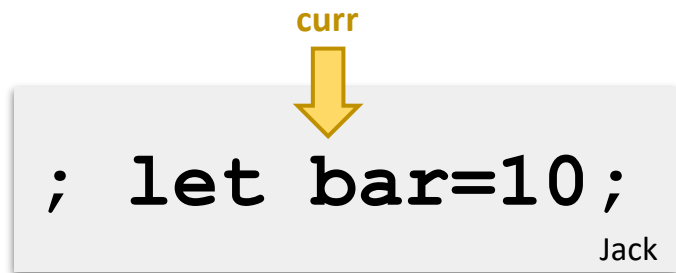
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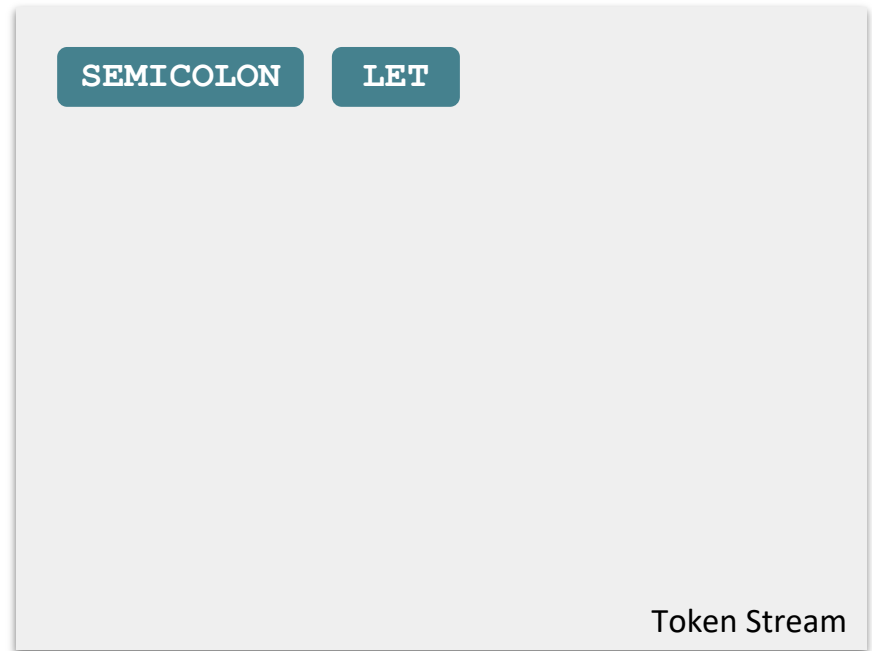


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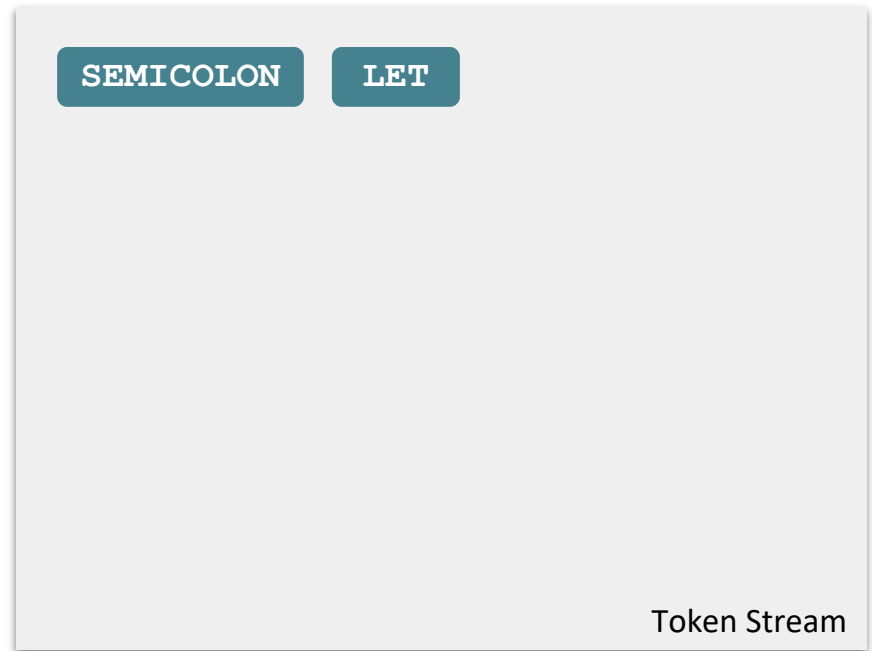
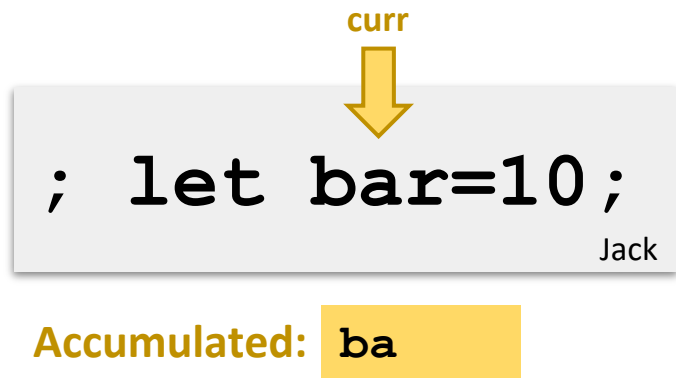


Accumulated: **b**



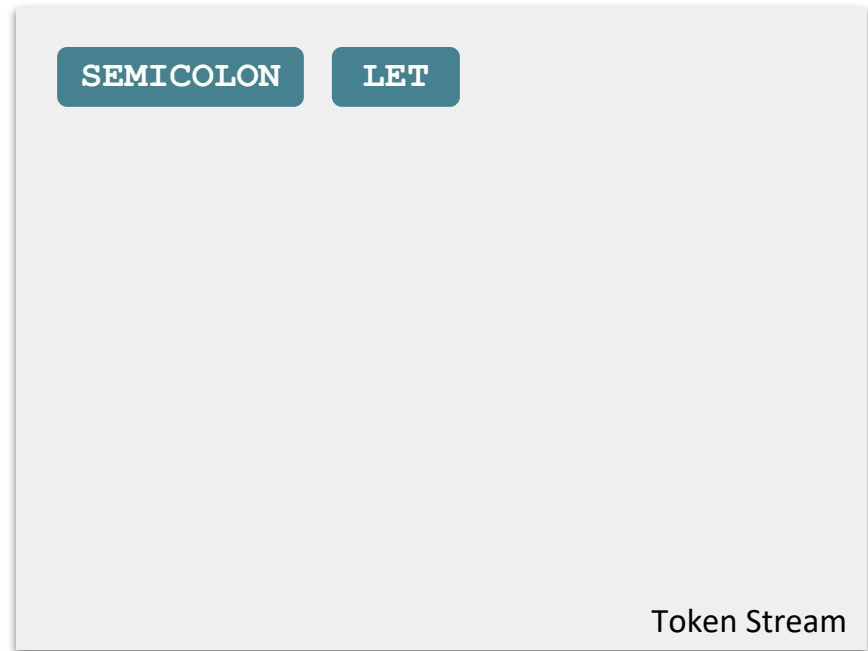
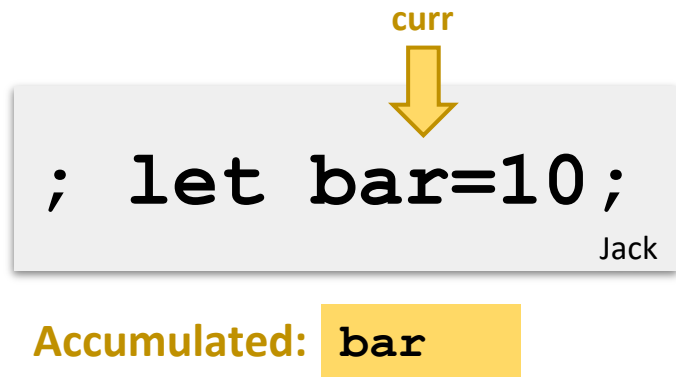
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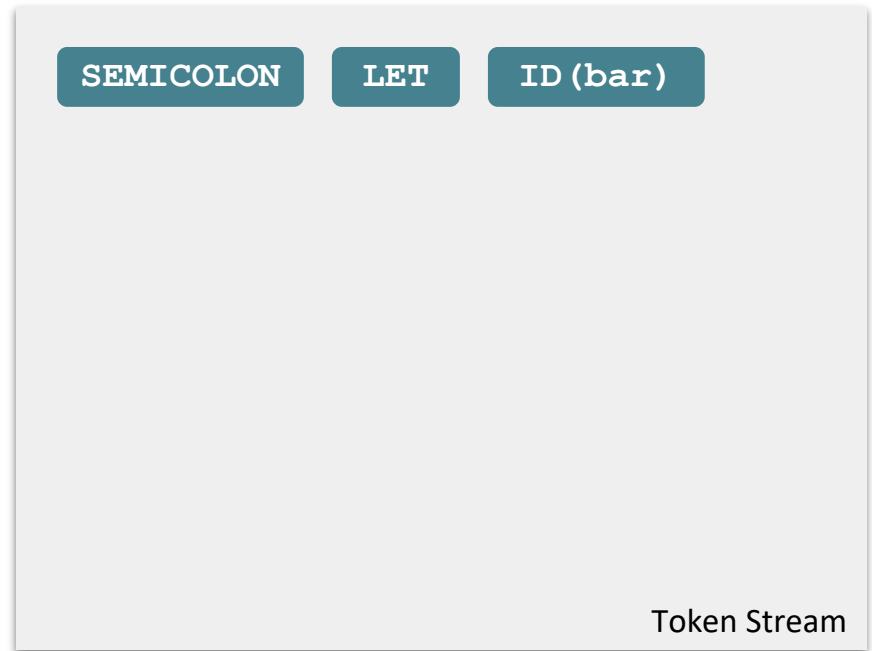
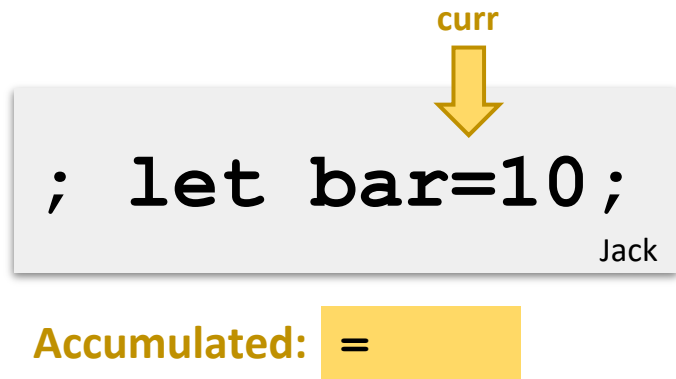
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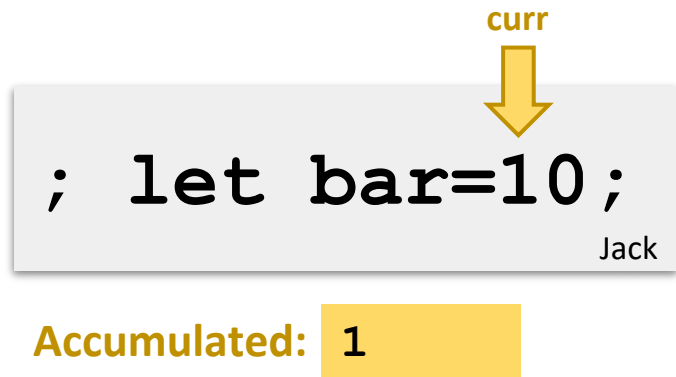
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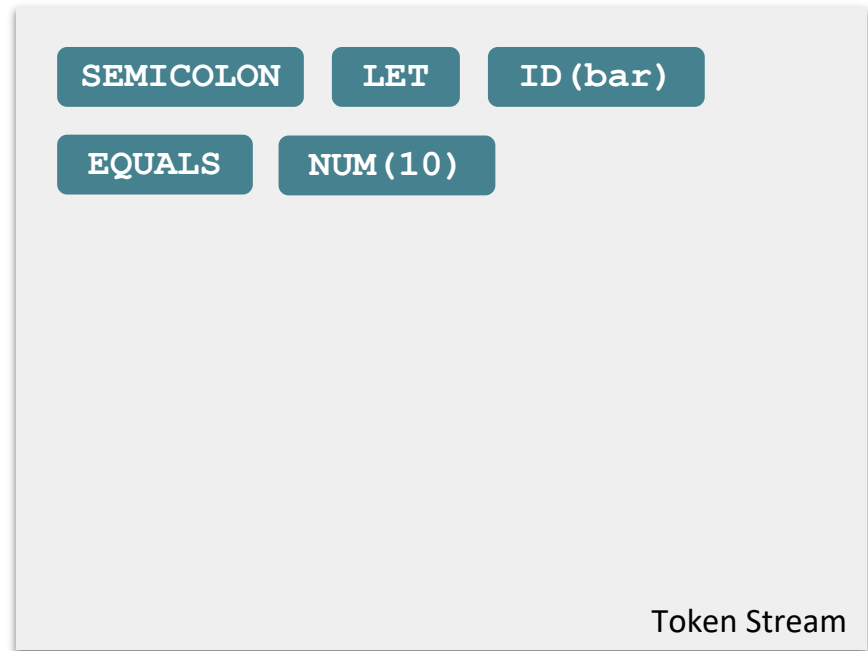
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The Scanner: How?

`; let bar=10;`
Jack

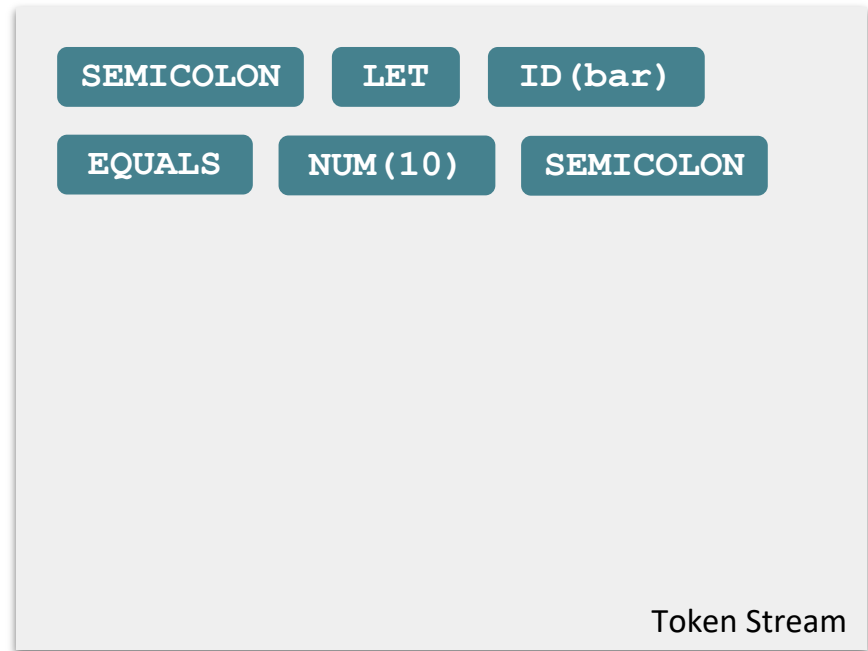
curr
↓

Accumulated: ;



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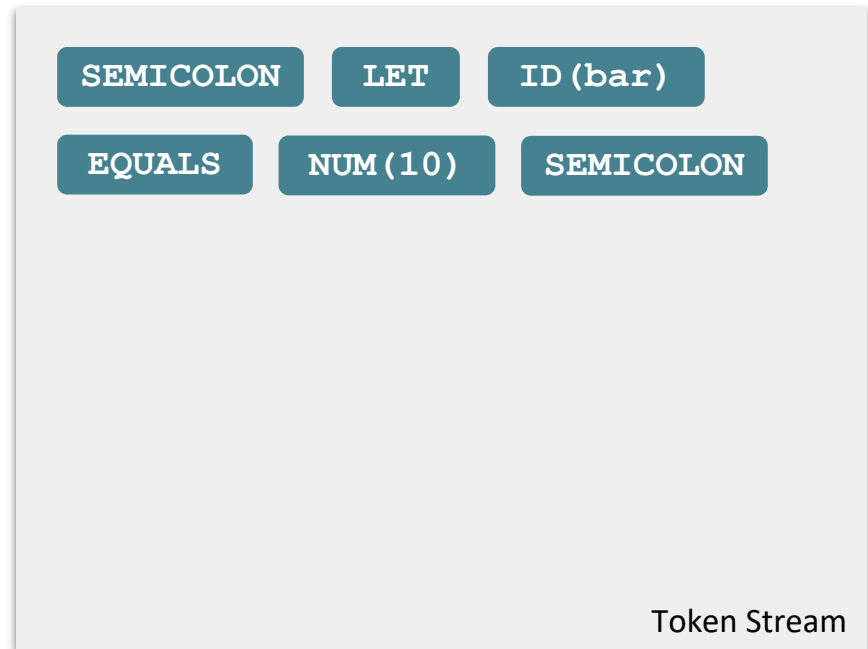
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The Scanner: How?

`; let bar=10;`
Jack

curr

Accumulated:



- ❖ How to distinguish built-in keywords (e.g., “let”) from identifiers (e.g., “bar”)?
 - When token is done, check against list of keywords

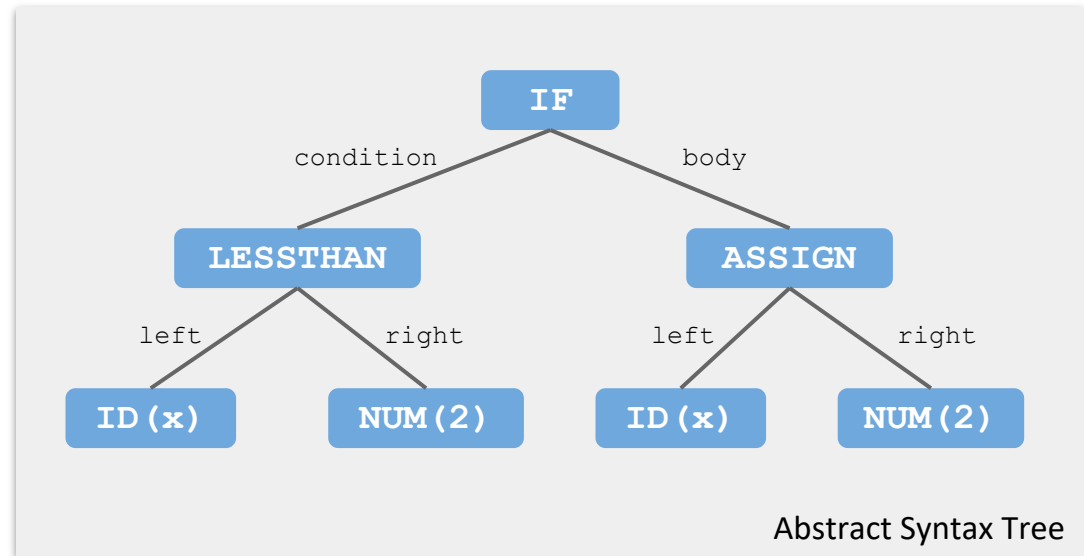
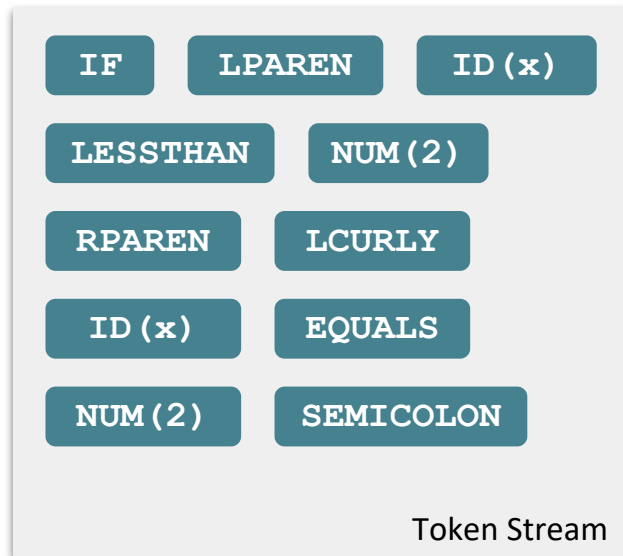
The Scanner: Why?

- ❖ Fundamentally: The compiler can't reason about a massive string, so we need to boil it down to its meaning
 - A great place to start is grouping characters that form a “word”
- ❖ Engineering-wise: Separation of concerns
 - A stream of tokens is an important abstraction for many file-processing tasks, not just compiling
 - Cleaning away whitespace and comments makes rest of compiler simpler

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- ❖ **Exploring the Compiler Phases**
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 - **Parser: Making Meaning From Tokens Through ASTs**
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The Parser



- ❖ Takes in the *flat* token stream and outputs a *structured* tree representation of program constructs
- ❖ Result: an **Abstract Syntax Tree**
 - Captures the structural features of the program
 - **Important distinction:** cares about **big-picture syntax** (E.g., entire `if` statement) rather than **nitty-gritty syntax** (E.g., semicolons, parentheses, even word “if” used to write that `if` statement)

Describing a Programming Language

- ❖ Many ways to define programming languages, some formal
 - We won't cover language definition in depth
 - See CSE 341, CSE 401, CSE 402
- ❖ Example: Statements vs. Expressions

Statements

Perform an action

- ❖ Assignment Statement

```
x = y;
```

- ❖ If Statement

```
if (x == 0) {  
    x = y;  
}
```

Expressions

Evaluate to a result

- ❖ Operators

```
x == 0;
```

- ❖ Variable

```
x
```

- ❖ Constant

```
24
```

Describing a Programming Language

- ❖ These broad categories lend themselves well to recursive definitions
 - Easily express all possible configurations of the language constructs

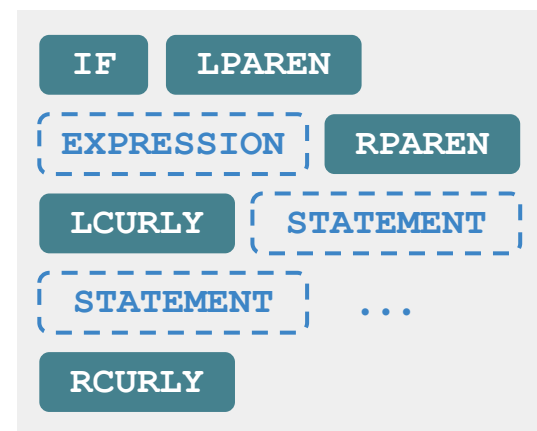
Symbolic Example

```
if (x == 0) {  
    x = y;  
}
```

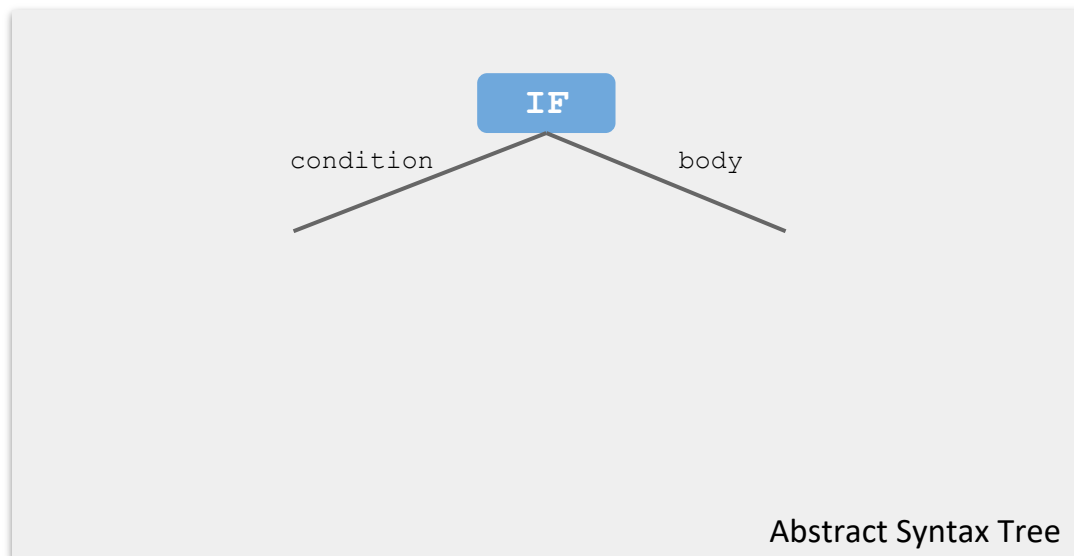
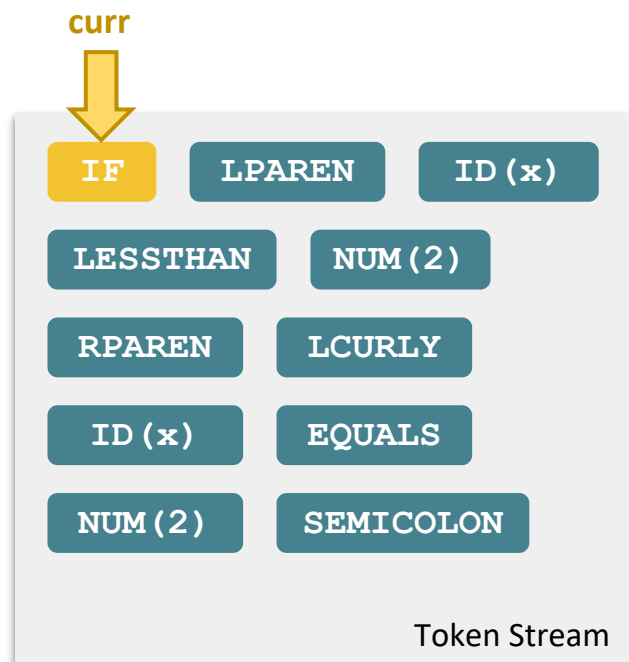
General Definition of an if Statement

```
if ( [ EXPRESSION ] )  
{  
    [ STATEMENT ]  
    [ STATEMENT ]  
    ...  
}
```

Token Stream Definition

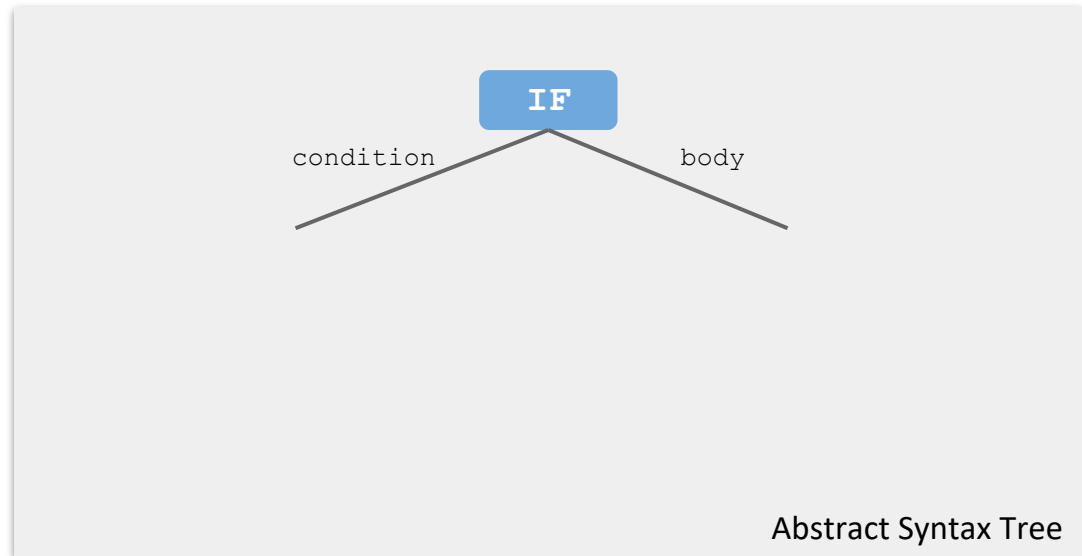
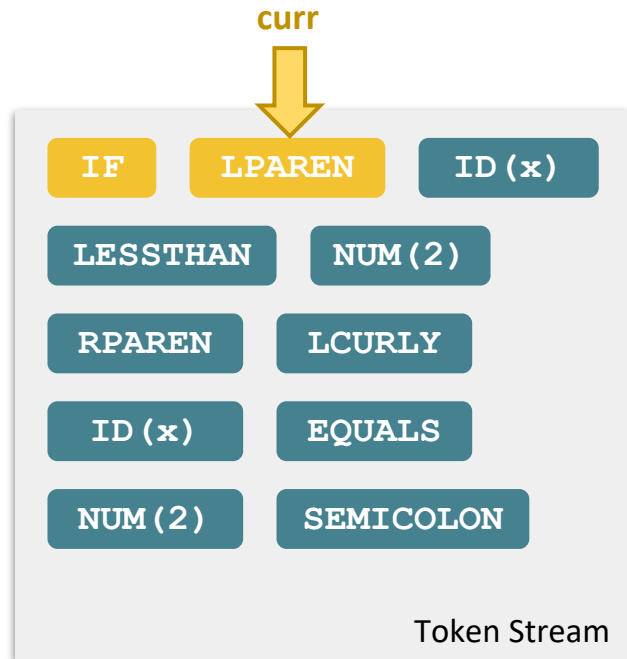


The Parser: How?



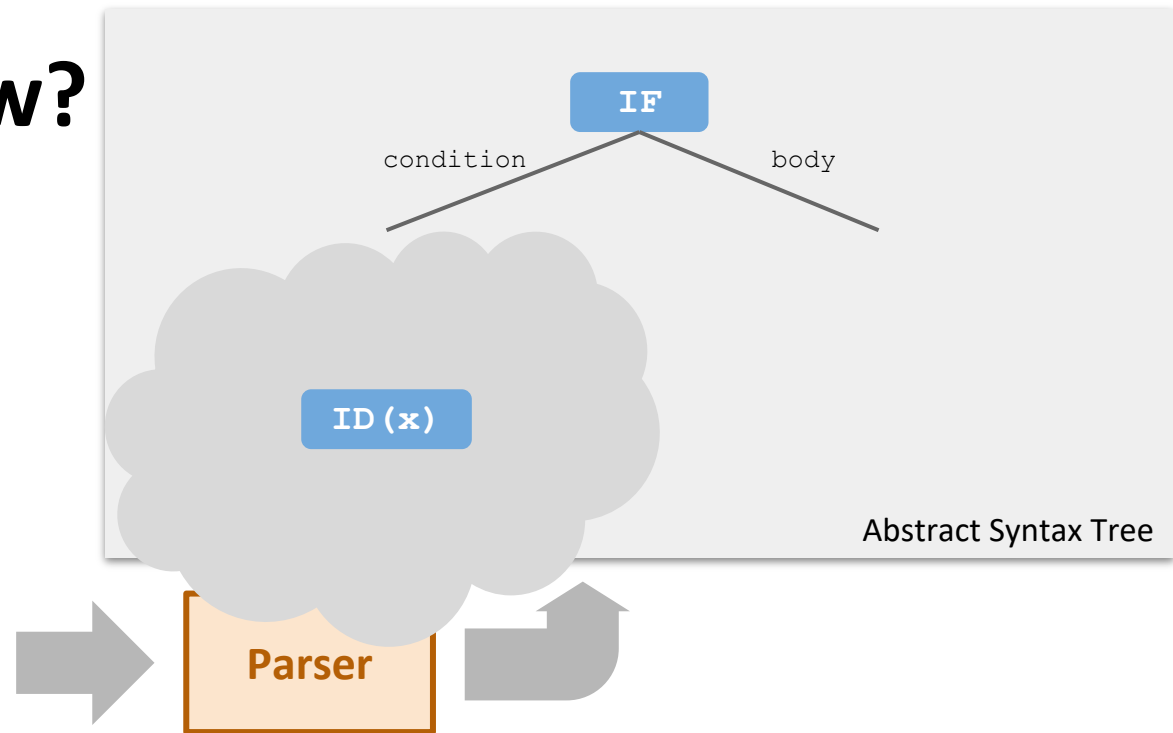
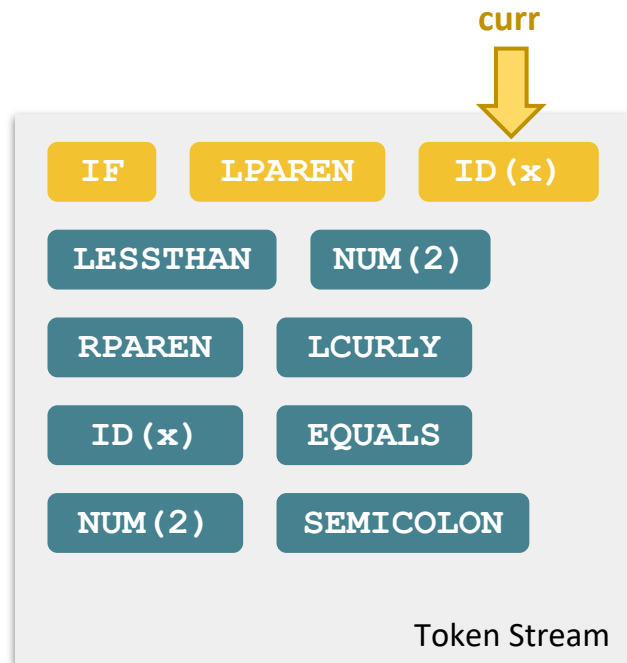
- ❖ Like a scanner: pass token stream, building up as we go
- ❖ Intuition: If we see **IF** and **LPAREN** , we are entering an if statement and next we must see a complete expression
 - Keep reading until we have a complete expression (recursively parse that) and attach on the condition side of the **IF**

The Parser: How?



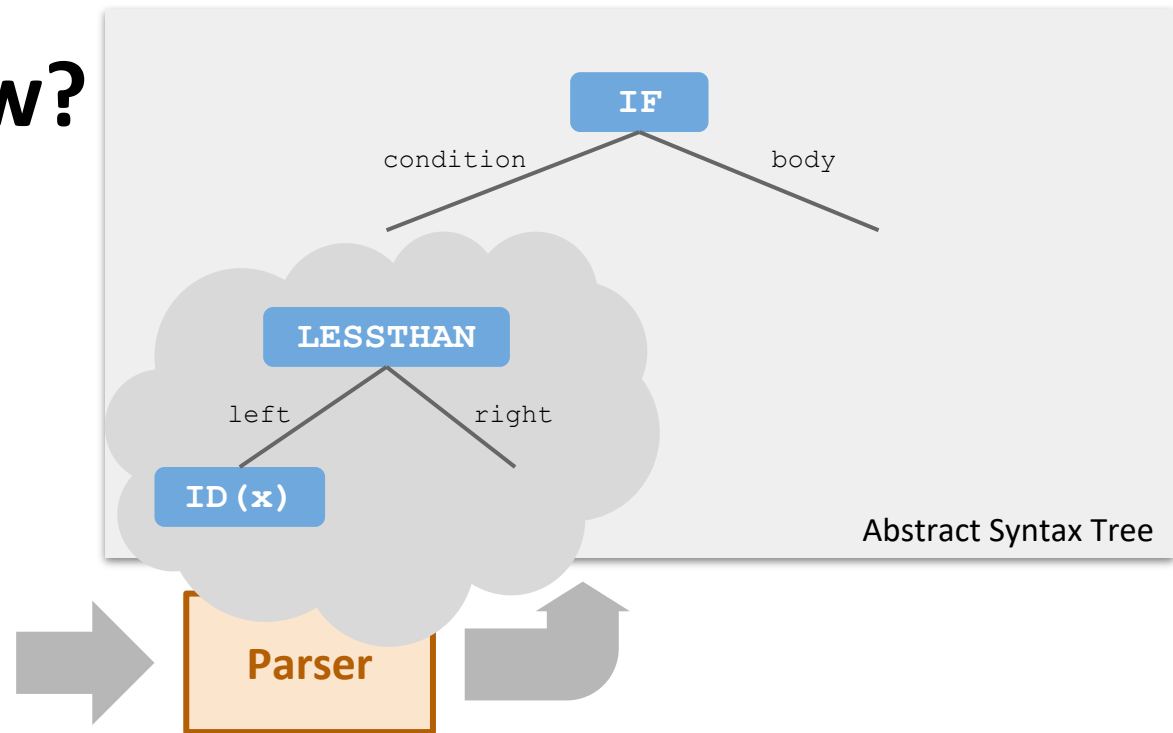
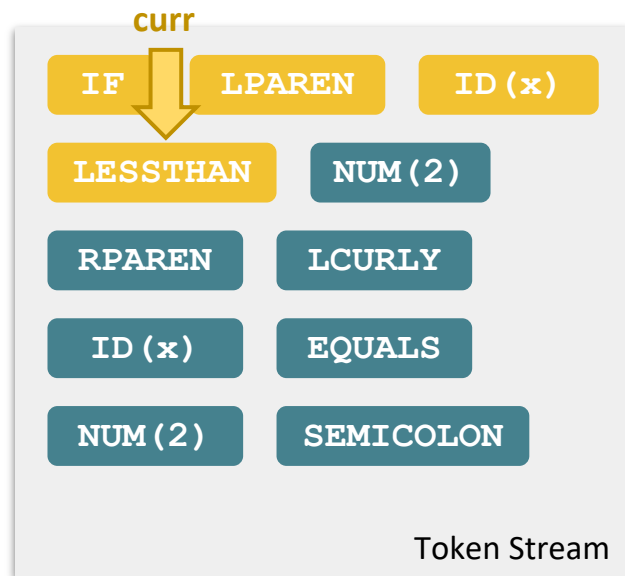
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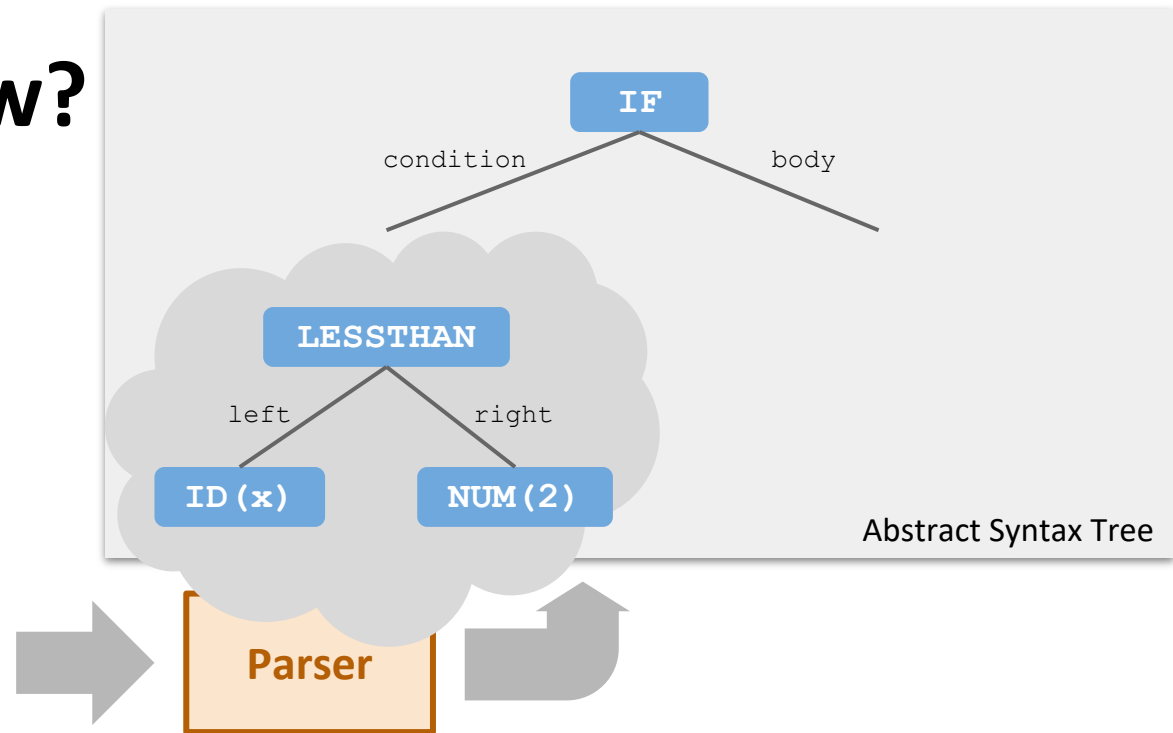
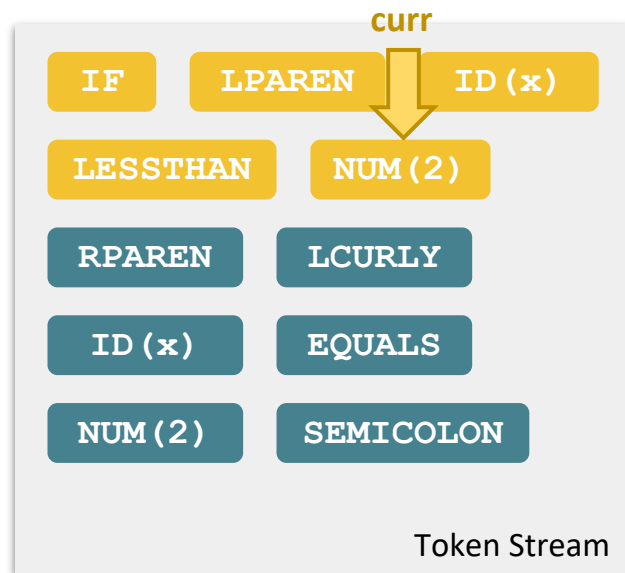
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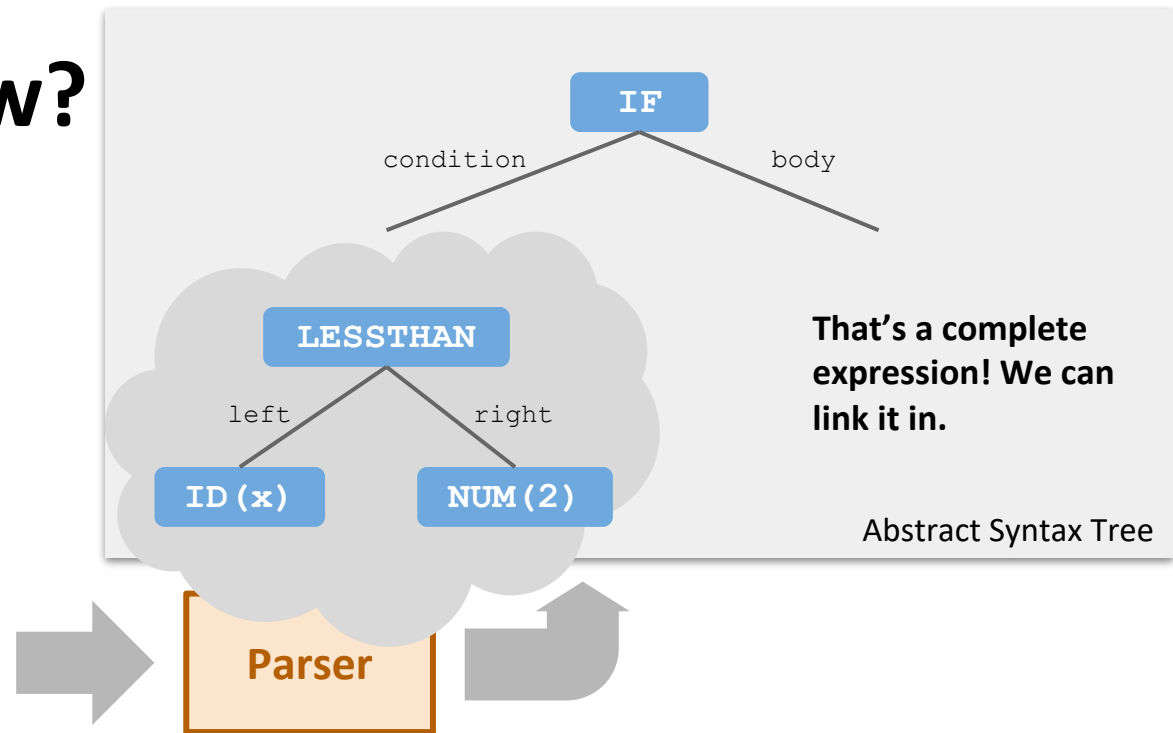
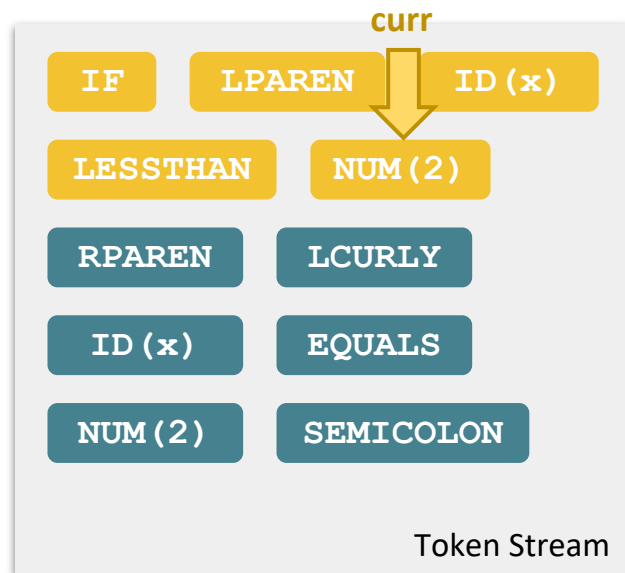
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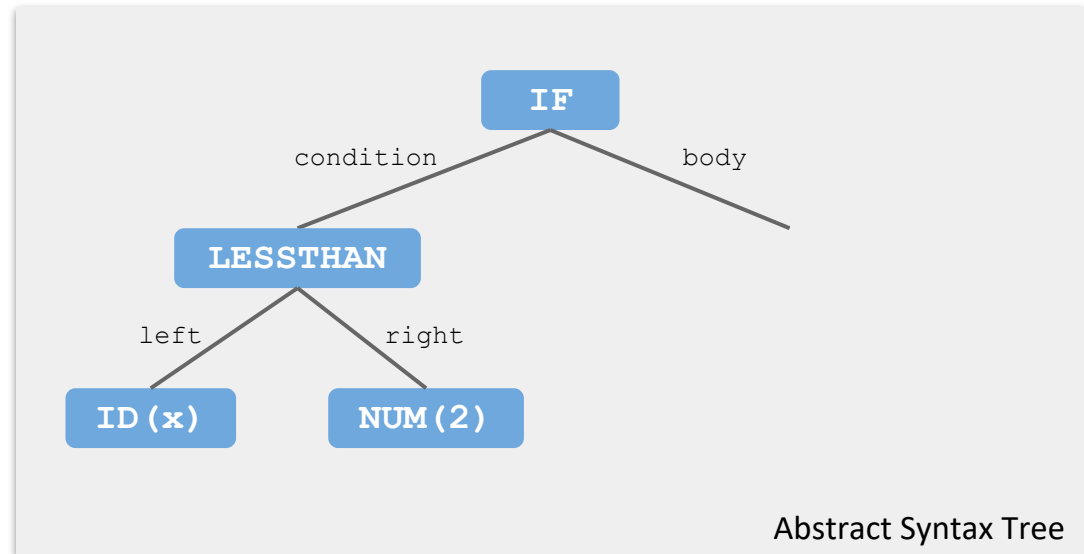
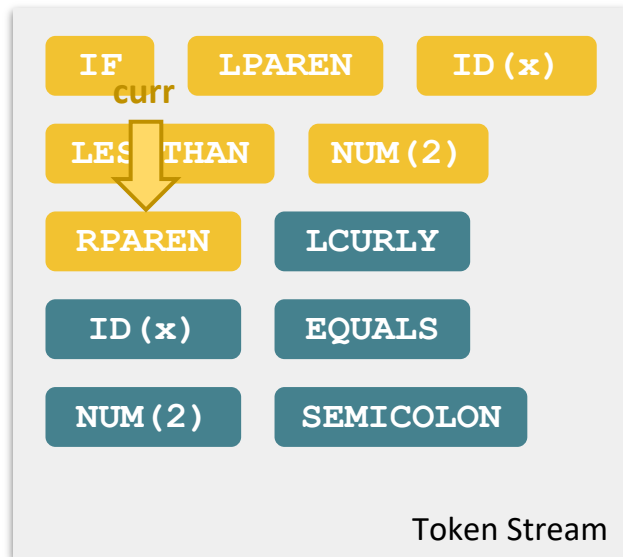
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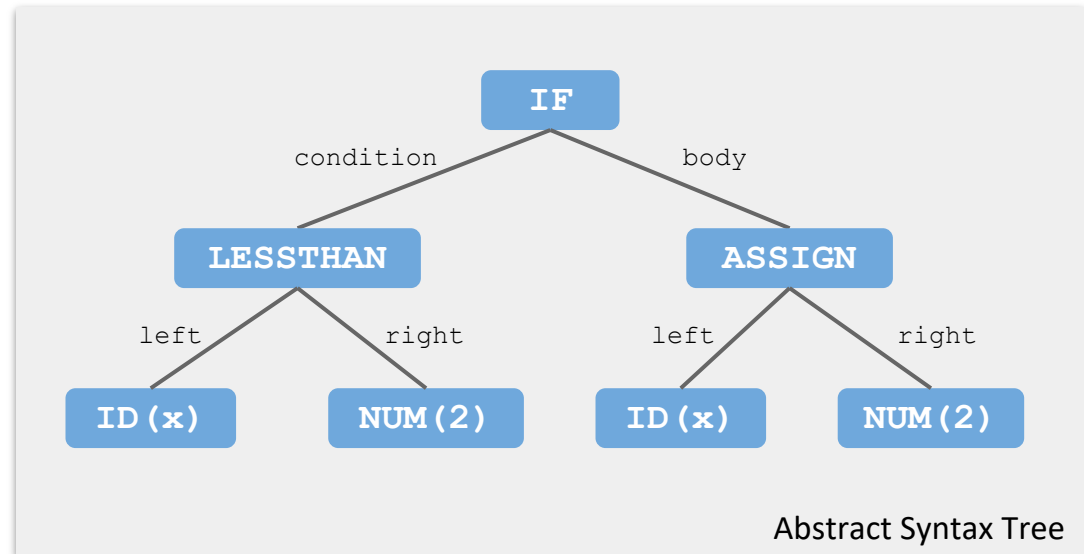
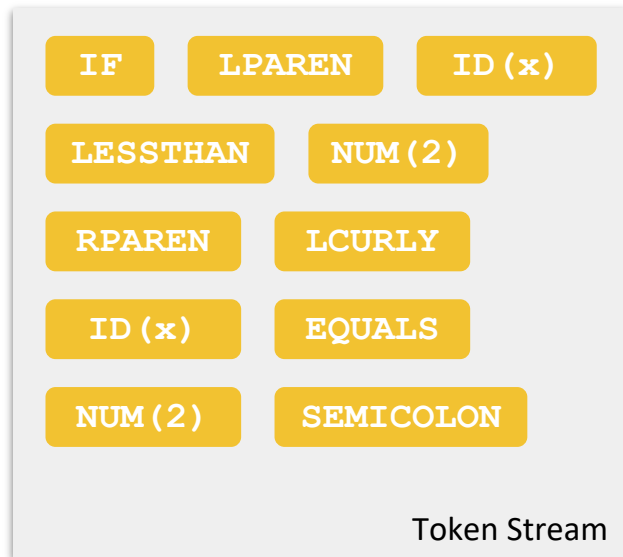
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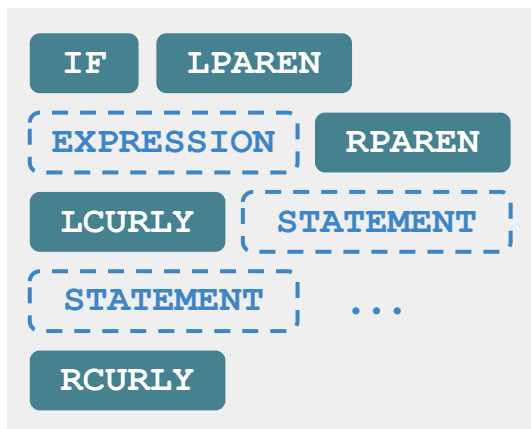


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The Parser: How?

- ❖ Implementing the Parser is essentially encoding the token stream definition, which can be recursive

Token Stream Definition



```

parseStatement () {
    ...
    if (currToken () == IF) {
        next () //consume "if"
        next () //consume "("

        // consumes tokens in expr
        e = parseExpression ()

        next () // consume ")"
        next () // consume "{"

        // consumes tokens in stmt
        s = parseStatement ()
        ...
        return new If (e, s)
    }
    ...
}

```

Lecture Outline

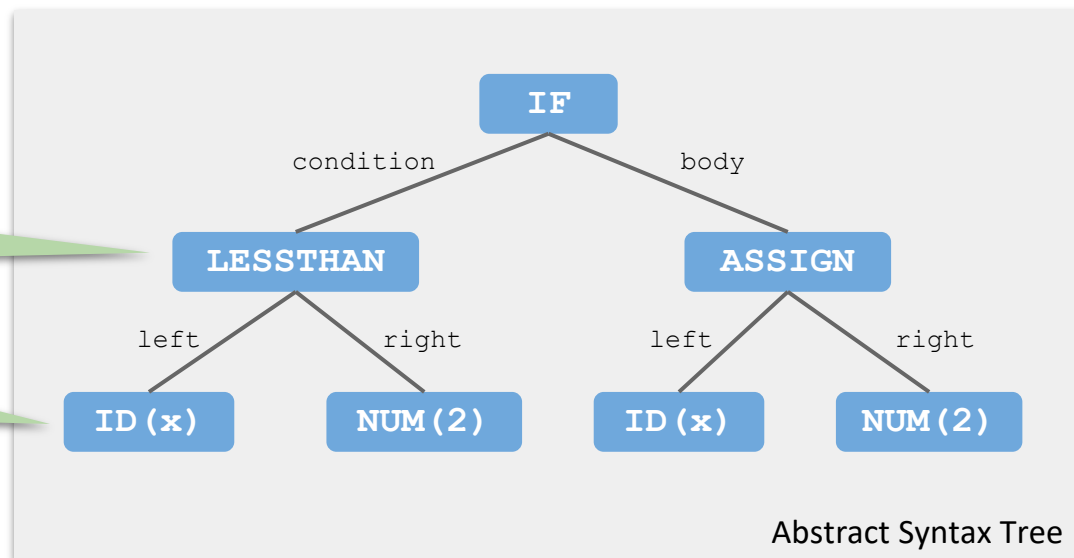
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 - Midterm Corrections, Professor Meeting Report

Type Checking (Semantic Analysis)

- ❖ Given the abstract syntax tree, run checks over it to ensure that it fits within constraints of the language
 - Do the types match up?
- ❖ Collect additional info for code generation, such as number and the type of arguments in each function

Does this expression evaluate to a Boolean?

Is the variable "x" defined at this point?



Optimization

- ❖ Code improvement: change correct code into semantically equivalent but “better” code
- ❖ Example: If something is computed every iteration of a while loop, the compiler could yank that computation out and compute it just once before entering the loop
 - Here, “better” means faster
- ❖ But requires caution: what if the value changes on each iteration of the loop?
 - “Semantically equivalent” means user sees same outcome

Code Generation

- ❖ One way to think of compiler is converting from string in source language to \rightarrow its actual, abstract “meaning”
- ❖ Code generation is converting that “meaning” into a string in the destination language
- ❖ At its core, all that the code generation phase does is read through the Abstract Syntax Tree and print a set of statements depending on the AST node

Lecture Outline

- ❖ Professional Networking in College
 - Benefits of Building Connections, Networking Strategies
- ❖ Introduction to Compilers
 - Scanner: Process of Tokenizing an Input File
 - Parser: Making Meaning From Tokens Through ASTs
 - Type Checking, Optimization, and Code Generation
- ❖ **Project 7 Overview**
 - **Midterm Corrections, Professor Meeting Report**

Project 7 Overview

❖ Part I: Midterm Corrections

- Due next Thursday (2/23) at 11:59pm (**no late days** can be used on midterm corrections)
- Open-notes, open-tools
- Only need to redo the problems that you missed
- 50% of the points you earn back from midterm corrections will be added to your original midterm score
- You can calculate your new midterm score using this formula:

$$\text{Original Midterm Score} + \frac{\text{New Midterm Score} - \text{Original Midterm Score}}{2}$$

❖ Part II: Professor Meeting Report

- Due in two weeks on 3/2 at 11:59pm
- Schedule the meeting as early as possible

Project 7, Part I: Midterm Corrections

- ❖ Review feedback from the course staff, celebrate the questions you got right, reflect on which areas you can continue to grow in
- ❖ If you think a problem was graded incorrectly, feel free to submit a regrade request on Gradescope
 - Don't be afraid to challenge our grading
 - This is a great learning opportunity for us all
- ❖ You can earn up to 50% of the points back that you missed on the midterm

Professor Meeting Report Discussion

In groups, spend 4-6 minutes discussing these questions:

- ❖ Which professors are you thinking about reaching out to? Why do you choose them?
- ❖ What questions would you ask to your professor? Why did you choose those questions?
- ❖ How can you apply the skill of meeting with professors in different contexts to help you succeed as a UW student? In your career?

Post-Lecture 14 Reminders

- ❖ **Project 6: Mock Exam Problem & Building a Computer due tonight (2/16) at 11:59pm**
- ❖ Project 7, Part I: Midterm Corrections due next Thursday (2/23) at 11:59pm
- ❖ Eric has office hours after class in CSE2 153
 - Feel free to post your questions on the Ed board as well